



*BY Developers FOR Developers*

**Storage Developer Conference**  
**September 22-23, 2020**

# **A Data Storage Diode for Classified Sites**

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# Kalray at SDC20

Kalray is well represented this year at SDC with 4 sessions! Please have a look.

- **A NVMe-oF Storage Diode for Classified Data Storage**  
Jean-Baptiste Riaux, Sr Field Application Engineer
- **High-performance RoCE/TCP Solutions for End-to-end NVMe-oF Communication**  
Jean-François Marie, Chief Solution Architect
- **Next Generation Datacenters Require Composable Architecture Enablers and Programmable Intelligence**  
Jean-François Marie, Chief Solution Architect
- **Smart Storage Adapter for Composable Architectures**  
Rémy Gauguey, Sr Software Architect



# **Abstract**

# Abstract

- Developing a “Storage Diode” by combining specific pieces of storage technologies such as HDF5, multipathing, ACL, user authentication (Kerberos, LDAP...) while leveraging NVMe-oF, is very useful for classified sites requiring remote and secure replication on NVMe SSDs.
- The storage diode is a dedicated storage system with two isolated Read and Write path, with guarantee of the data integrity. Leveraging dual port NVMe drives and the parallelism of advanced processors, this paper reviews how to fully isolate channels at both logical and physical levels, and dedicate write-only and read-only path to storage devices over a NVMe-oF fabric.

This technique allows restricted/classified computing center to push (write) data to the storage diode, assuring the path to the outside world can be only be accessed in Read-Only.



# The Presenter

# About the Presenter



Jean-Baptiste Riaux is a Sr Field Application Engineer at Kalray for the Data Center Business Unit.

He specializes in storage and HPC applications, and has a strong experience in deploying NVMe and NVMe-oF storage solutions.

Before joining Kalray, Jean-Baptiste gained an extensive experience of the storage and HPC domains in DDN & Intel.



# **Context & Requirements**

# What are HPC Dark Sites ?

- HPC Dark sites: classified datacenters closed from outside world running HPC clusters.
- HPC cluster:
  - Compute cluster sharing a unique parallel filesystem (shared disk FS, distributed FS)
  - Dozens of IO nodes, thousands of clients
  - The core of storage is a parallel filesystem such as Lustre or GPFS (PB of data)
  - GPU farms and specialized nodes for AI workflows



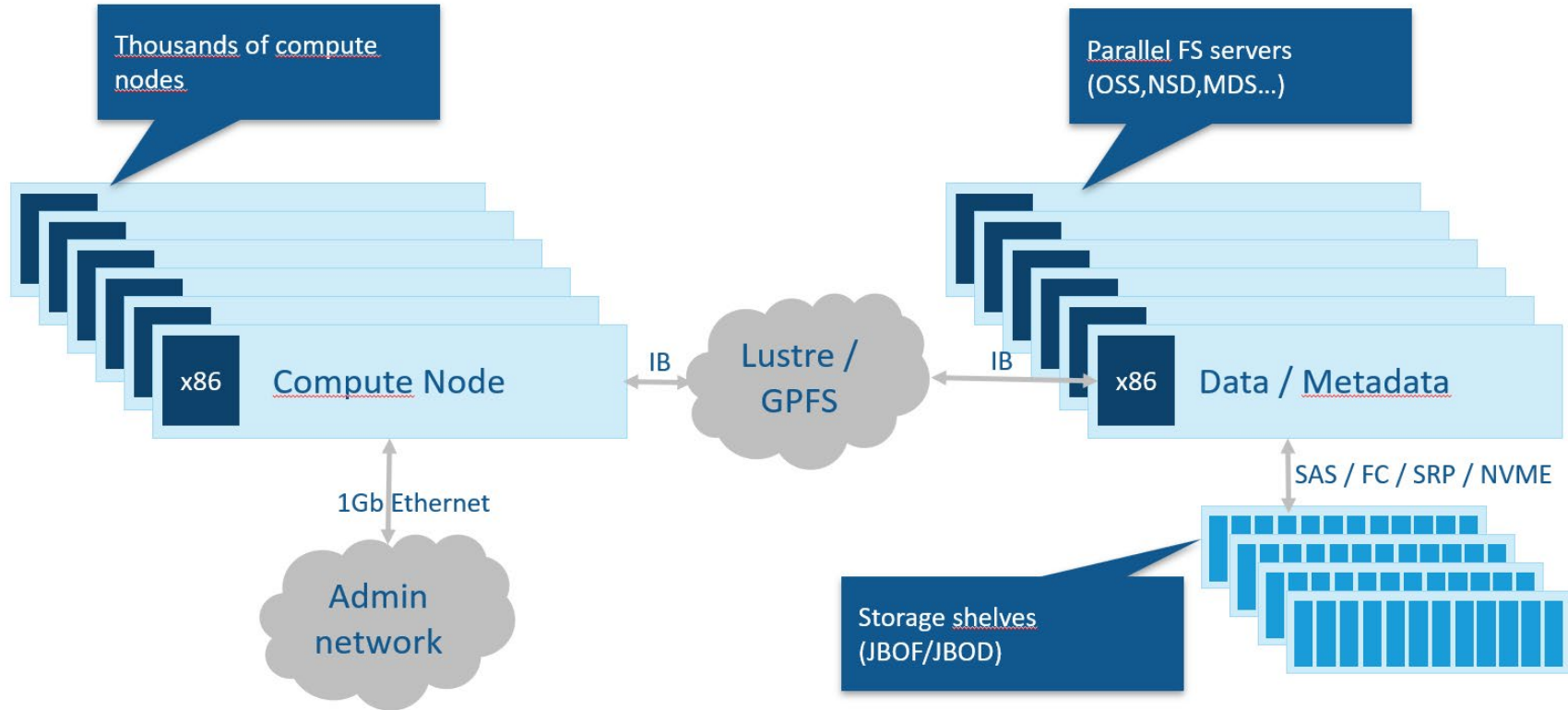
# HPC Dark Sites Challenges

- Dark sites can have high security concerns, such as:
  - No internet access
  - No VPN access (or for a very few restricted system engineers + selected HPC users / researchers)
- Generated data is not always for internal use only

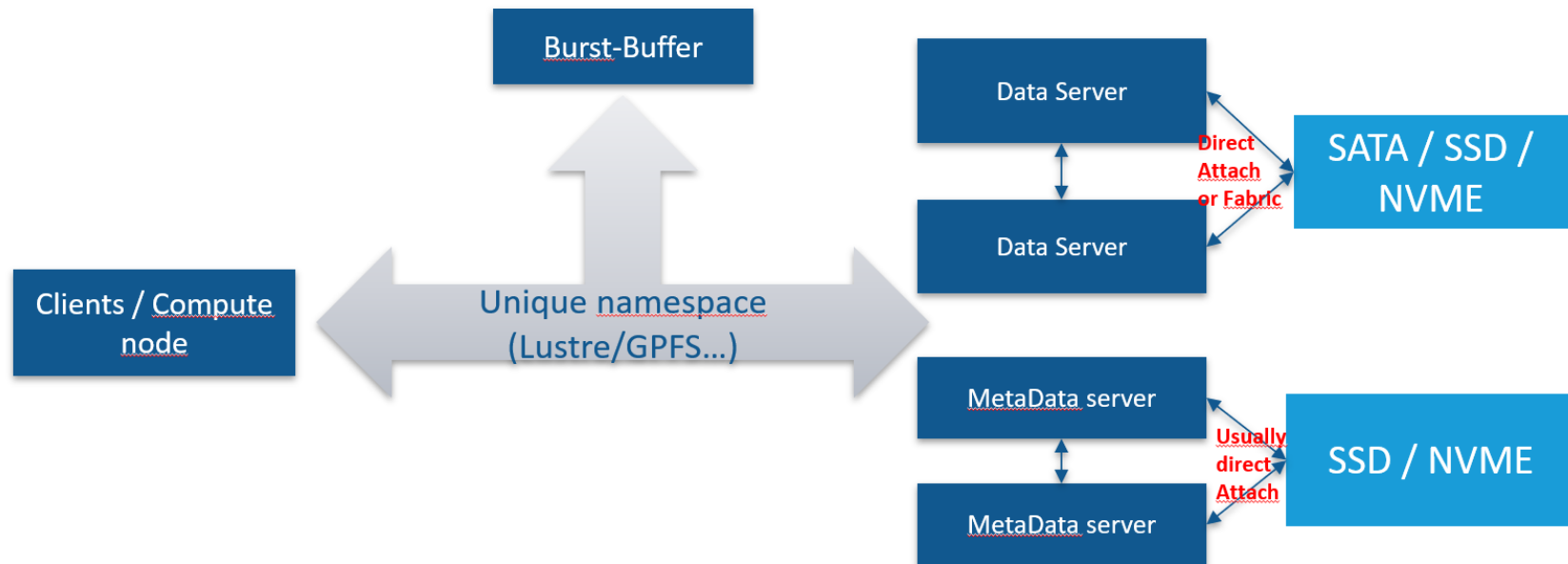
# Providing Data Safely to the Outside

- Controlled cold copy on drive(s)
- Special node(s) outside for remote copy using multiple layers of encryption
- Direct optical link to a close and secured datacenter

# What is a HPC cluster ?



# What is a HPC filesystem ?



# Limitations

- Limit as much as possible the accesses required for an outside copy
- Constraints:
  - Can't be interconnected to the parallel filesystem
  - Can't have access to PFS namespace
  - Can't be connected to the same admin network
  - Can't be a file based access protocol (network based)



Block-based access solution

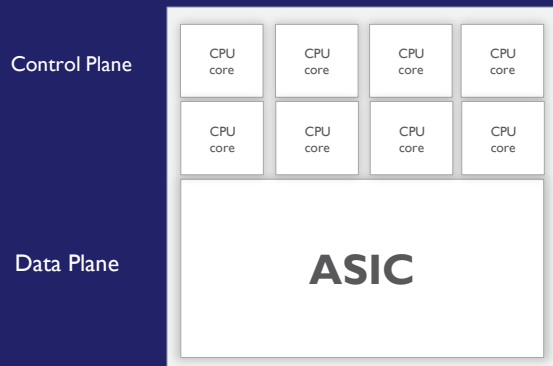


**How can Kalray Help?**

# COOLIDGE™: the ULTIMATE I/O Processor

Why Coolidge is a Revolution vs Competition ?

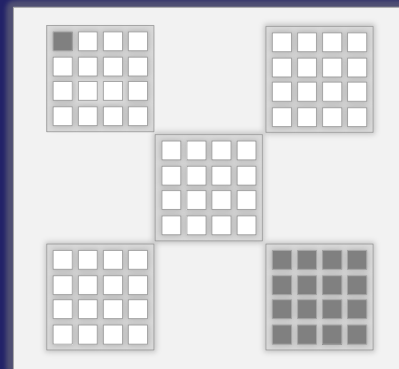
## “xPU” Usual Approach



**CONS**

- A few power hungry RISC CPU cores
- CPU flexibility limited to control plane
- Data plane is “hardwired” –  
No new services / no possible evolution!

## Kalray's MPPA®3 Coolidge™



80 highly efficient VLIW independent **CPU** cores, gathered into **5 clusters**, running at **1.2Ghz**, connected to high speed fabrics & high speed interfaces.

### </> Fully programmable

Control Plane / Mgt Plane – Linux – 16 cores  
Data Plane - 64 cores

**PROS**



**Power efficiency**

25W Typ



**High Speed I/O**

2x100Gbs,PCIGen4,DDR4



**Top Performance**

**Any workload**

200KDMIPs, 25TOPS



**Functional Isolation & Safety**

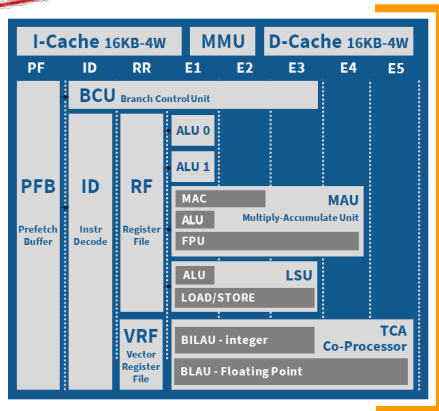
**Safety**

Secure Islands, Encrypt/Decrypt, Secure Boot

# MPPA® COOLIDGE Architecture

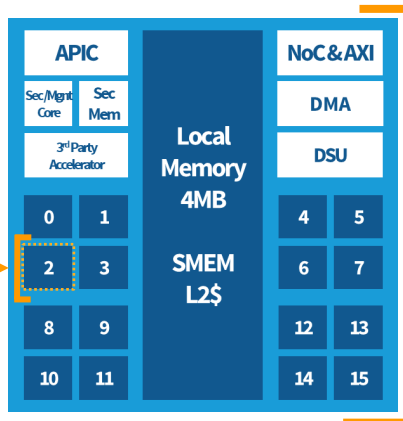
The I/O Processor for Next Gen Intelligent Systems

**PATENTED**



## 3<sup>RD</sup> GENERATION KALRAY CORE

- VLIW 64-bit core
- 6-issue VLIW architecture
- MMU + I&D cache (16KB+16KB)
- 16-bit/32-bit/64-bit IEEE 754-2008 FPU
- Vision/CNN Co-processor (TCA)



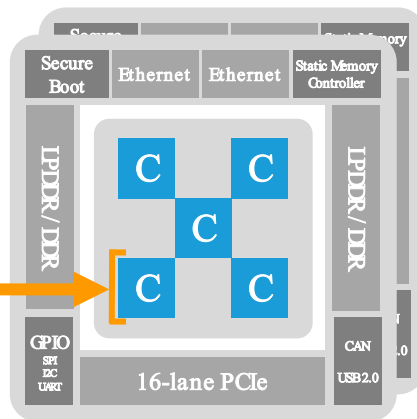
## CLUSTER

### Architecture

- 16 cores
- 1 safety/security dedicated core
- 600 to 1200 MHz

### Memory

- L1 cache coherency (configurable)
- 4MB configurable memory (L2 cache)
- 256 bits / bandwidth up to 614GB/s)



## MULTI CLUSTER ARCHITECTURE

### 5 Clusters: 80 cores + 80 co-processors

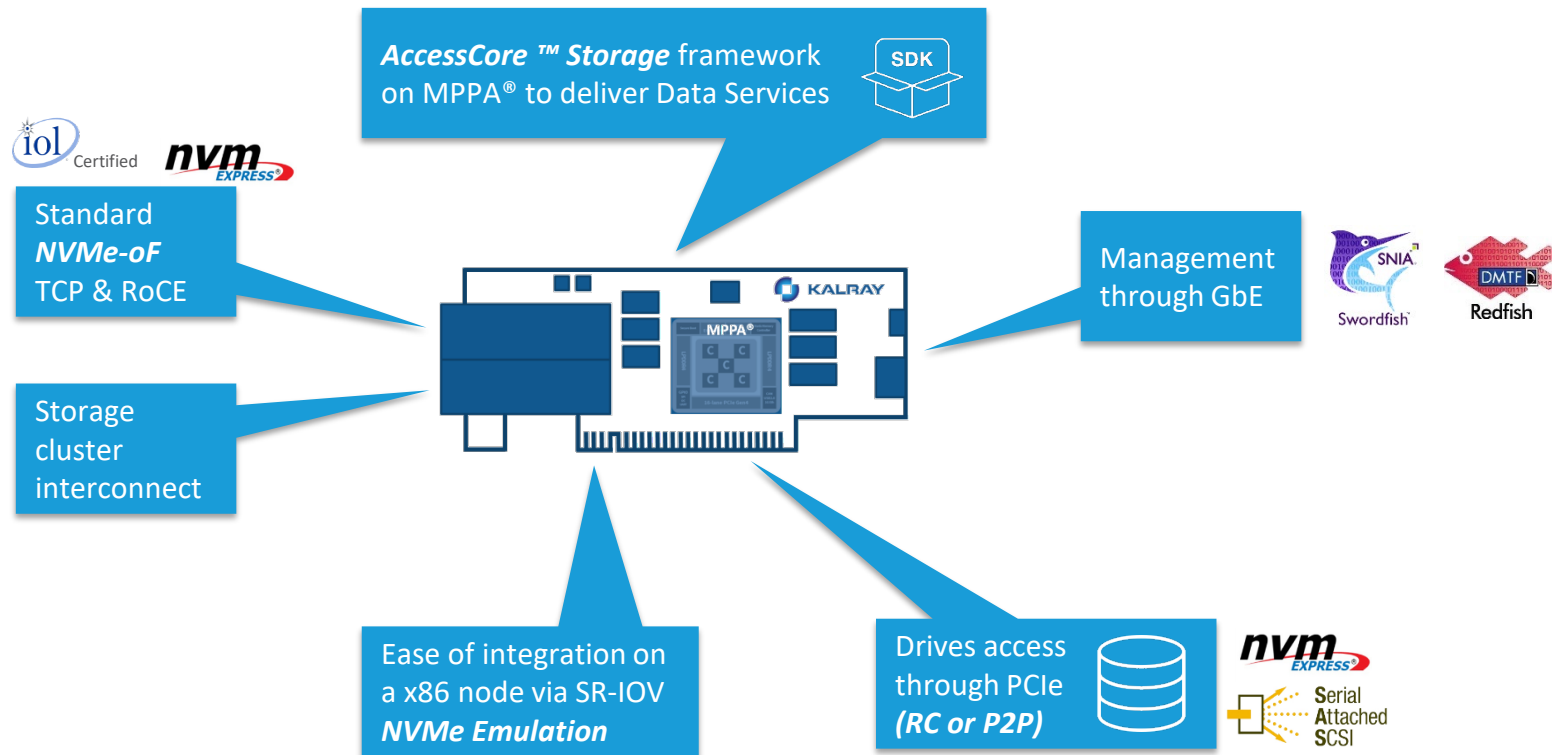
- Load Balancer / Packet Parser
- 2x100Gbps Ethernet
- PCIeGen4
- DDR4 - 3200

### AXI Bus + NoC Bus

- L2 refill in DDR and direct access to DDR from clusters
- DMA-based highly efficient data connection



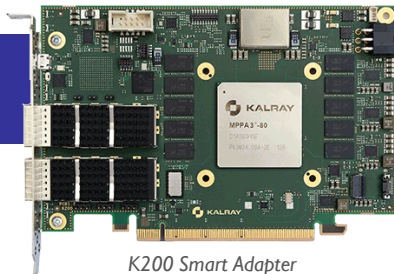
# Kalray Smart Storage Adapter



# Kalray Smart Storage Adapter Solution

## K200 & K200-LP

manufactured by **wistron**



K200 Smart Adapter

### 2 Form Factors

- FHHL (Full Height) - K200 - Single Slot
- HHHL (Low Profile) - K200-LP  
Single or Double Slots

### 2 Modes

- Stand-alone
- Host CPU co-processor  
/ "host-agnostic" support

### Manycore Architecture

- 80 VLIW cores @ 1.2 Ghz
- 5 Clusters x16 cores

### High Speed Ethernet

- 2x100GbE / 8x25 GbE

### Certified NVMe-oF Stack

- NVMe-oF 1.1 (Target, Initiator)
- RoCE v1/v2, TCP

### Advanced SSD interface

- PCIe-Gen4
- NVMe 1.1 to 1.4 SSDs  
No need for CMB
- Dual port SSD support

### Agnostic Host Support

- NVMe Driver

### DDR-3200

- 8GB to 32GB

### H/W Accelerators

- Encryption / Decryption
- Hashing (SHA-256, SHA-3)
- Erasure Coding

### Low Power

- 35W (single slot)
- 65W (double slot)

### Key figures per card

- Random R/W RoCE: **4-6 MIOPS**
- Random R/W TCP: **2-4 MIOPS**
- Sequential R/W (RoCE&TCP):  
**25GB/s**
- Latency (RoCE/TCP): **10 /30  
usec**



**AccessCore®**  
Open Software & Tools

### Open Software Environment

- Linux / SPDK Control Plane (16 Cores)
- Fully Programmable Data Plane (64 Cores)
- Storage, Network and Compute Services  
(AI,DSP,NVMe,NVMe-oF,ROCE,TCP, RAID, de-dup,...)

### Agnostic Host Support

- NVMe Driver

### + Extra compute available

- @ 3MIOPS, 50% cores available !
- Storage Services (RAID, de-dedup ...)
- AI
- Analytics ...

# AccessCore® for Storage & Networking

## PROGRAMMABILITY

- Full programmability on data, control & management planes
  - Control & Management plane : Linux (typical : 1 Cluster - 16 cores)
  - Data plane : Cluster OS (light POSIX OS) (typical: 1 to 4 Clusters – 16 to 64 cores)

## EFFICIENCY

- Run to completion full dataplane
  - From network functions to NVMe stack on light OS cores
- True inline processing
  - No need for x86 pre/post processing

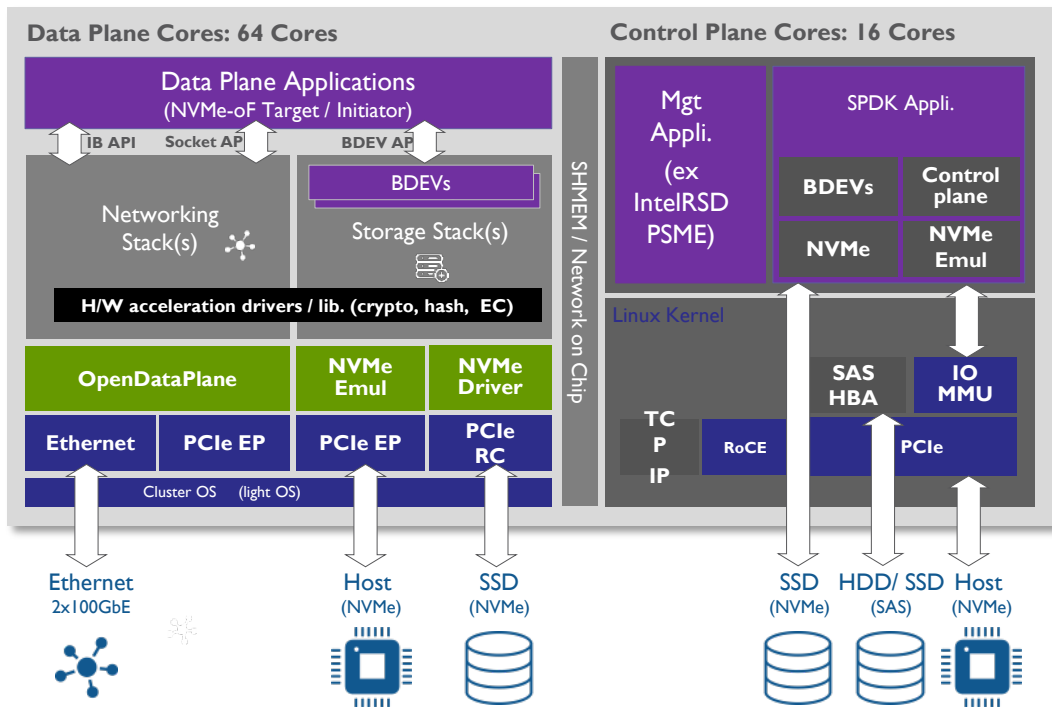
## STANDARDIZED

- Hardware interfaces
  - NVMe emulation
- Software APIs & tool chain
  - Linux APIs: SPDK, virtio, ibverbs ...
  - Data plane APIs: sockets, SPDK nvme lib, SPDK BDEV, ODP
  - Libraries : ISA-L, Buildroot, binutils



# AccessCore®

## A fully flexible software environment



- A complete & modular software framework
- Based on an optimized SPDK for both data plane **AND** control plane
- Open to partners



# Example of NVMe-oF (RoCE/TCP) JBOF

## Hyper Optimized JBOF (no x86)

- JBOF Chassis :
  - Stand-alone
  - 2U – 1200W Redundant
  - 24 U.2 NVMe SSDs
  - 6xPCIe Gen3 x16
- Kalray Smart Controller Cards
  - 2 to 6 Cards
- BMC chip – AST2500 (ASpeed)
- 1Gbps management interface

## Lymma JBOF Reference Platform White Label NVMe-oF (RoCE/TCP) JBOF



wistron

NVMe SSDs

Redundant Power

System Cooling FANs

PCIe Card Cages 12



# Toward a true & efficient composable disaggregated Infrastructure

## HIGHER PERFORMANCE

- Leverage Kalray cards performance and exploit full NVMe SSD capabilities
- Offload x86 from heavy storage stacks

## LOWER COST

- Switch to a true **C**omposable **D**isaggregated **I**nfrastructure with commodity components
- Optimize HCI nodes efficiency

## FULLY FLEXIBLE

- Fully programmable data plane
- Data Plane additional storage services based on SPDK framework (EC, caching...)

## FUTURE PROOF

- Leverage standard NVMe-oF protocols
- Compliant with other NVMe-oF appliances
- Ease of in-the-field update



# Global Design Overview

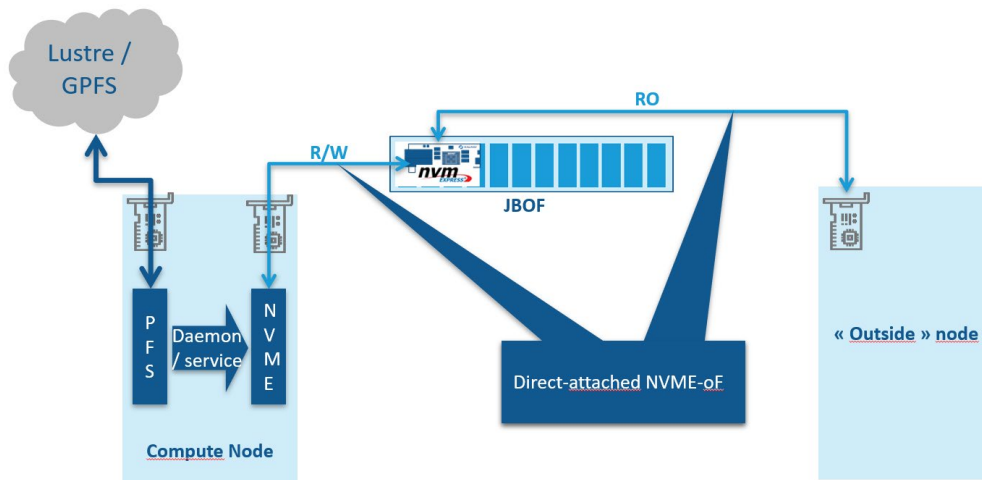
# Design overview

## Dark-site node

Lustre/GPFS client with a second high-speed interface (40/100G)

The second interface is used for NVME-oF

Periodically writes data from PFS to NVME devices



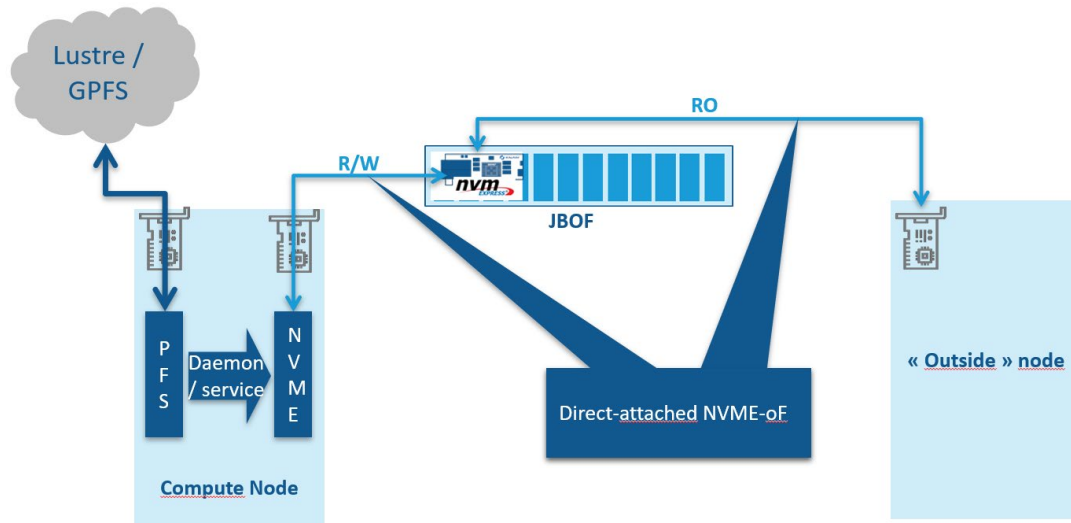
## Outside node

Single high speed interface used for NVME-oF  
Direct-attached NVMe-oF for security



# Design overview

- Looks easy but two problems remain:
  - What about inside the FBOF ?
  - With block level access only, still need to ensure file locking at filesystem level.
- Inside the FBOF: separation at PCI level on dual ported drives
  - Using PCI switches for either flow separation or resiliency

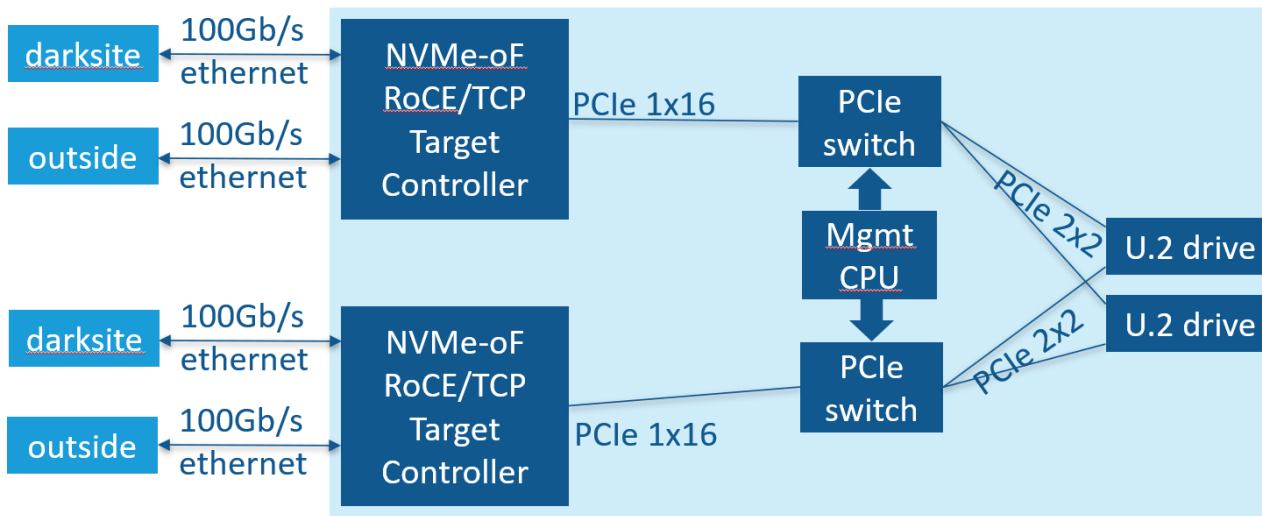




# FBOF Parameters

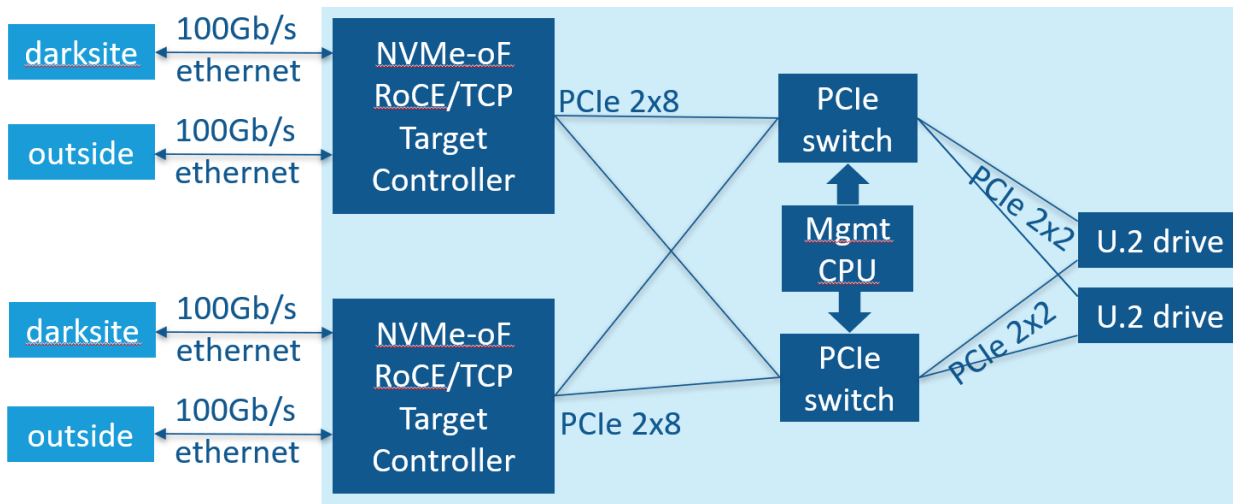
# FBOF PCI Configuration

- U.2 dual ported drives to ensure flow separation between the two physical paths
- At PCI level, multiple configurations/usages can be done depending on resiliency vs isolation



# Custom security with K200

- Target controller PCI bifurcation can be also defined between 2x8, 4x4 and 1x16 PCI lanes
- Using our own board (Kalray's K200) allows us to add custom security



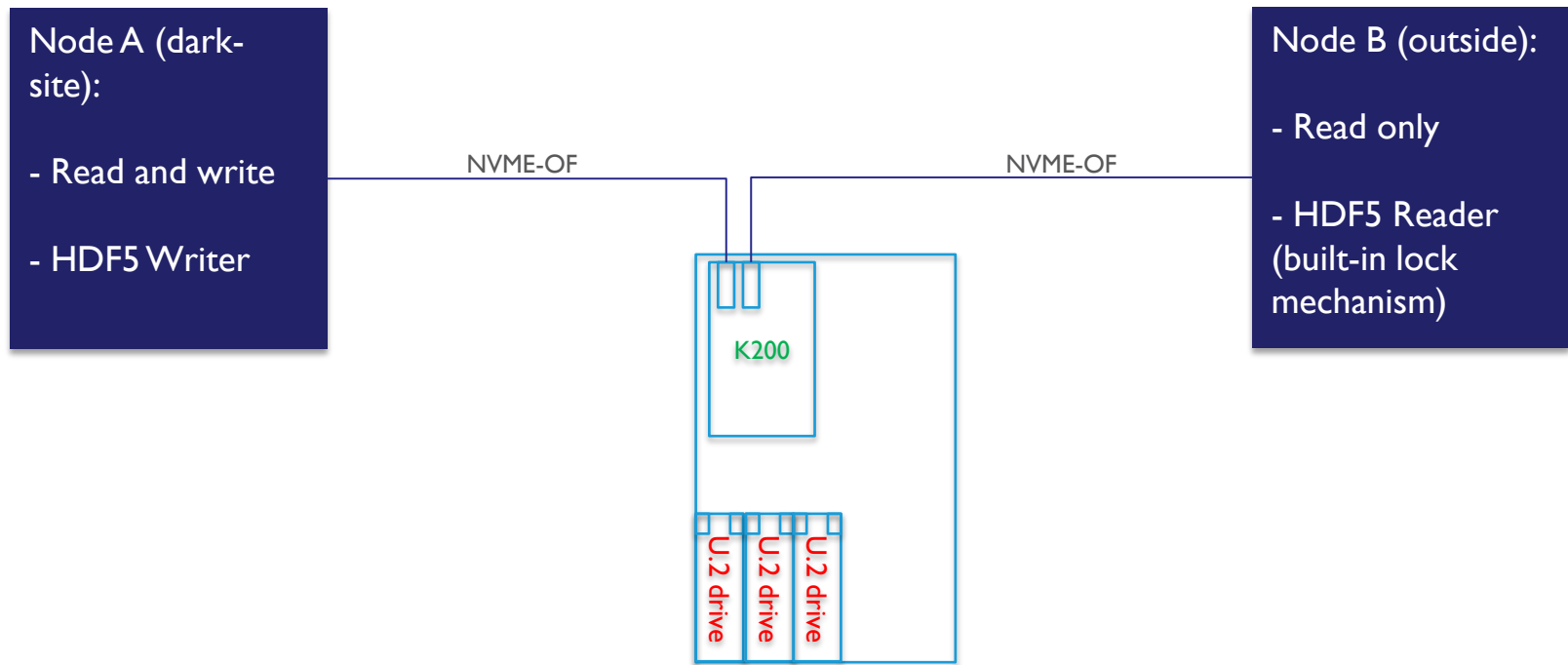


# **The Provided Solution**

# Design Options

- Two nodes connected to a NVME target controller, in direct connect
  - One node as read and write
  - One node as read only
  - Guaranteed separated data paths
  - Block level protocol (NVME)
- 
- Need for a filesystem using internal locking at file level or containers
    - HDF5 has been chosen

# Architecture View



# Key Elements

- Ideal for WORM workflows (Write Once Read Many)
- Benchmarking is necessary to find a good balance between:
  - Writers and number of drives (x86 saturation vs nvme saturation)
  - Drive IOPS in write and read (specialized read drives or not)
  - Number of readers to a single target





# **Key Benefits**

# Key benefits

## SECURITY FOCUSED

- Guaranteed read-only access from outside world
- PFS namespace hidden
- Restricted NVME access

## SCALABLE

- Easy to scale, possibility to start small and add:
  - K200 cards
  - U.2 drives
  - Writers nodes
  - Reader nodes
  - Full diode

## FUTURE PROOF

- As MPPA® processor is fully programmable
- Ease to update
- Add custom code (security / feature)



**Q&A on slack**  
**Thank you!**